

Towards an Algebra of Routing Tables

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Australian Government

Department of Broadband, Communications
and the Digital Economy

Australian Research Council

NICTA Members



Department of State and
Regional Development



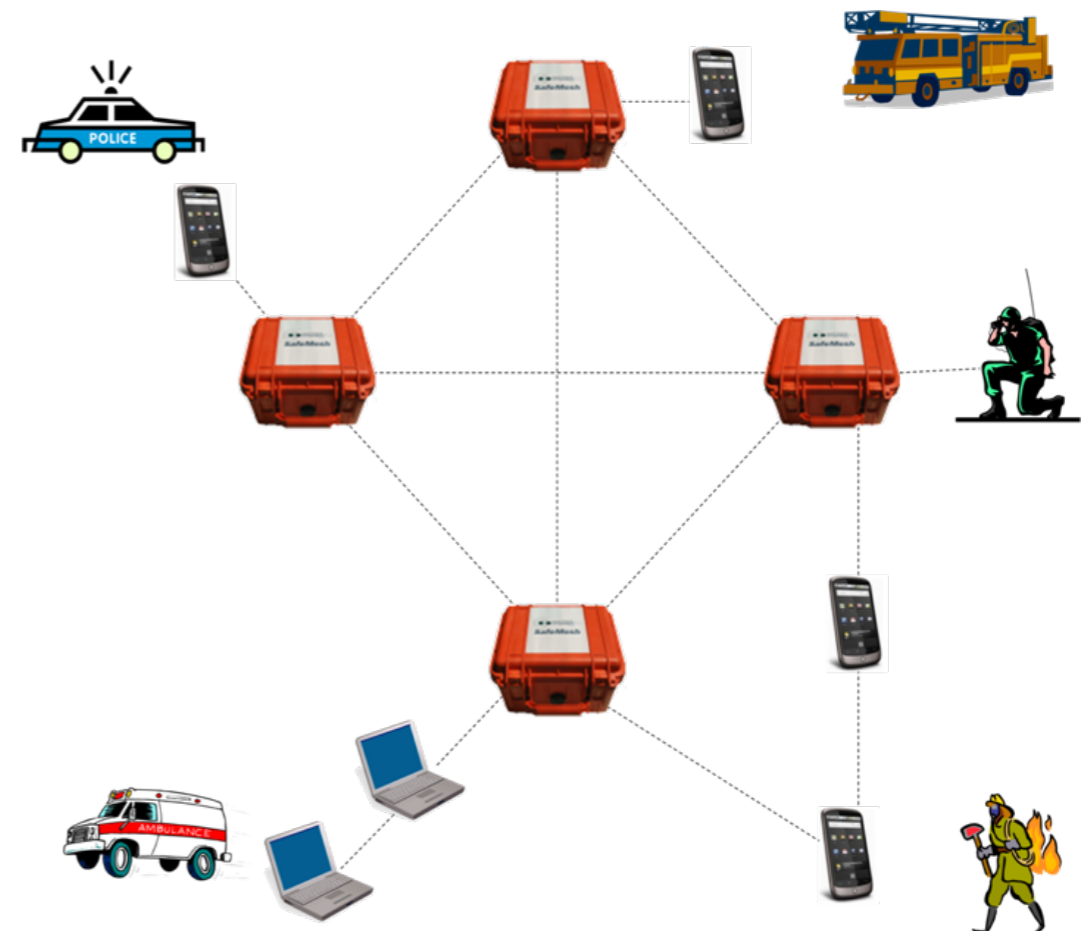
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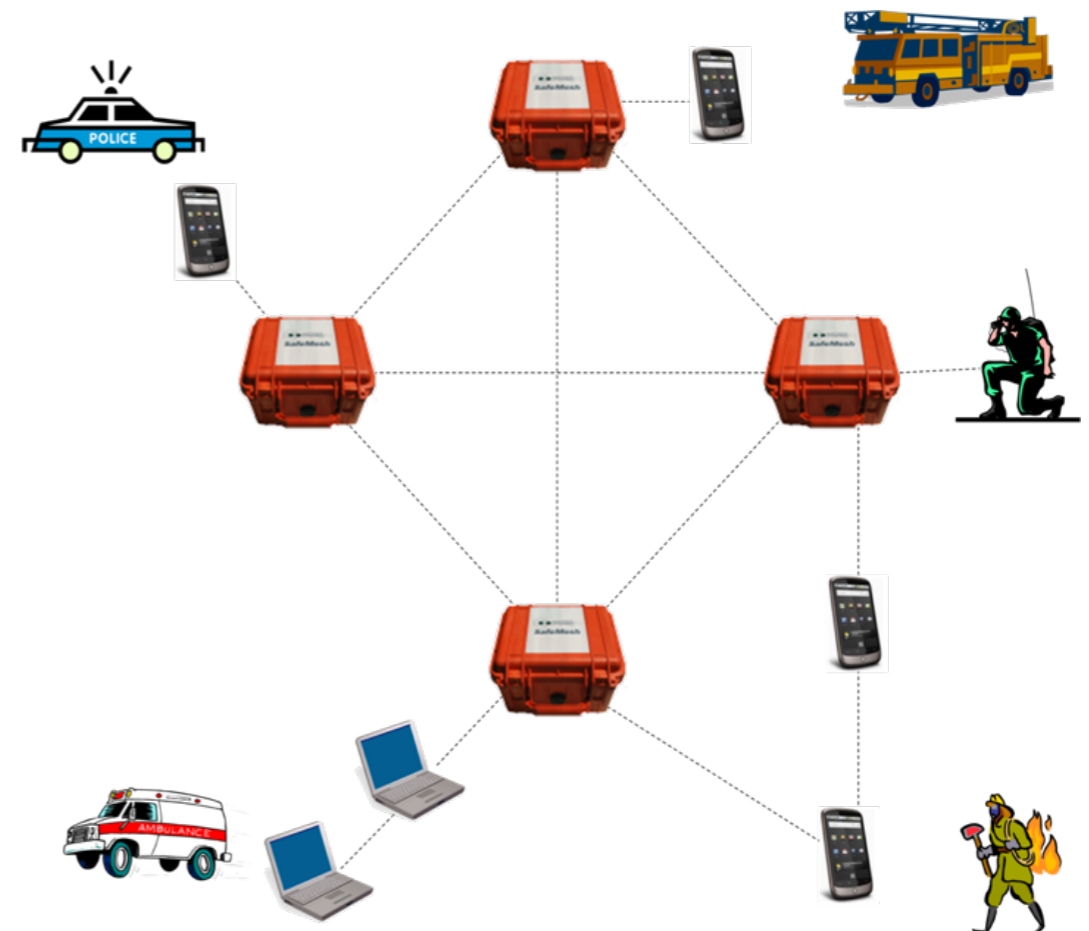
What is the Problem?

- **Wireless Mesh Networks**
 - key advantage: no backhaul wiring required
 - quick and low cost deployment
- **Applications**
 - public safety (e.g. CCTV)
 - emergencies (e.g. earthquakes)
 - mobile phone services
 - transportation
 - mining
 - military actions/counter terrorism
 - ...



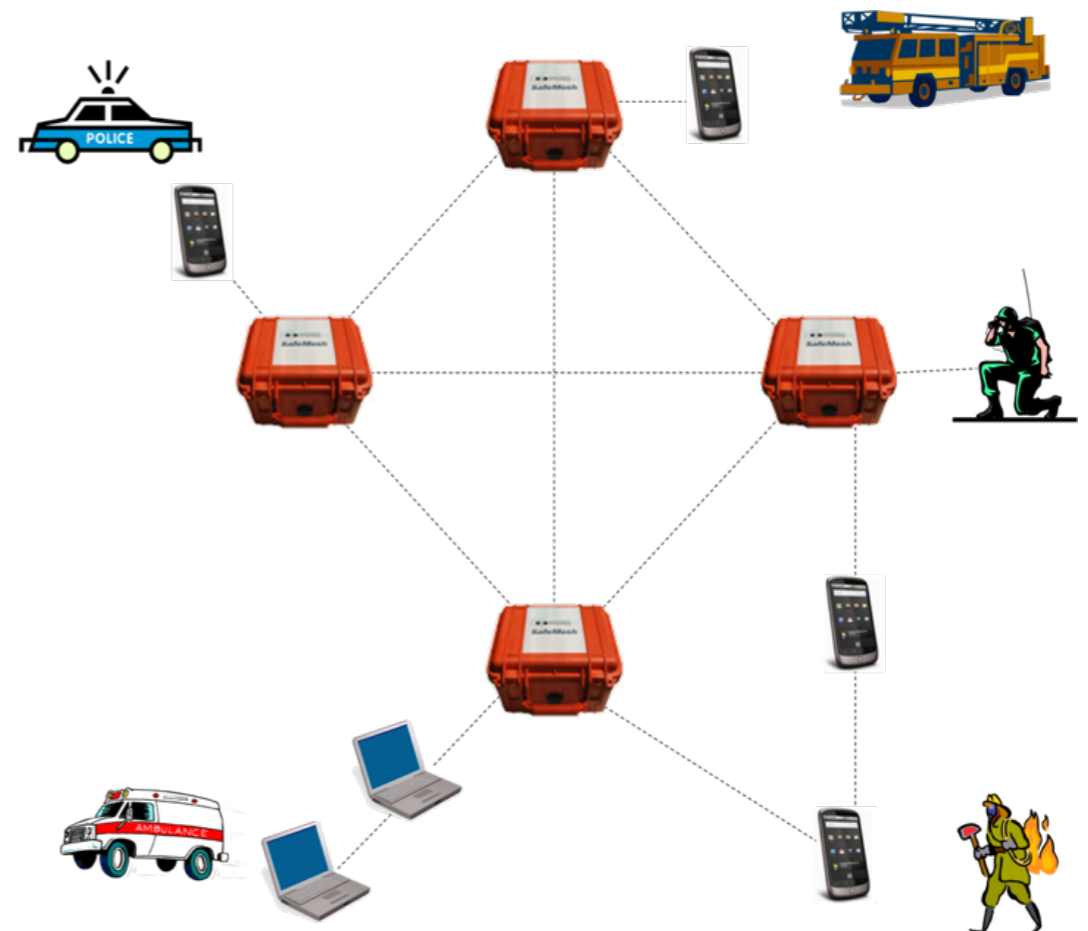
What is the Problem?

- WMNs promise to be fully
 - self-configuring
 - self-healing
 - self-optimising



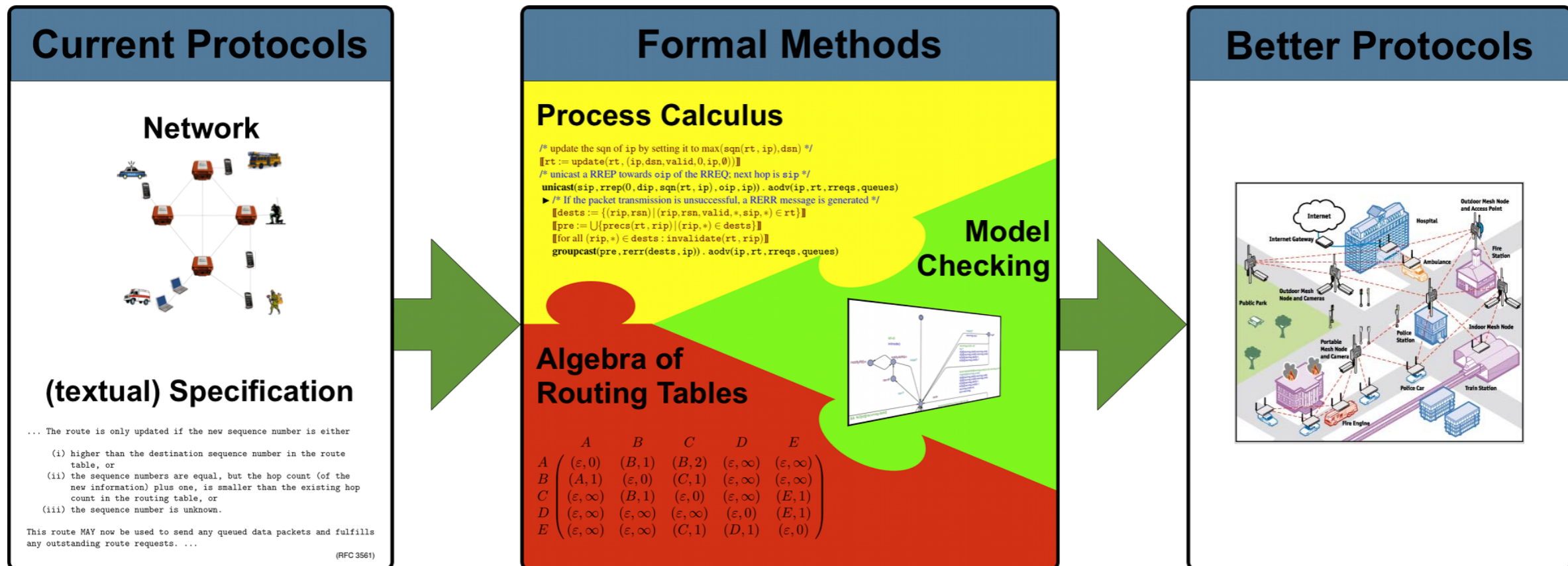
What is the Problem?

- WMNs promise to be fully
 - self-configuring
 - self-healing
 - self-optimising
- **THAT IS NOT TRUE**
(in reality)
- Limitations in reliability and performance
- Limitations confirmed by
 - end users (e.g. police)
 - own experiments
 - Cisco, Motorola, Firetide, ...
 - industry



- **Goal**
 - model, analyse, verify and increase the performance of wireless mesh protocols
 - develop suitable formal methods techniques
- **Benefits**
 - more reliable protocols
 - finding and fixing bugs
 - better performance
 - proving correctness
 - reduce “time-to-market”

- Main Methods used so far
 - process algebra
 - model checking
 - **algebra of routing tables**



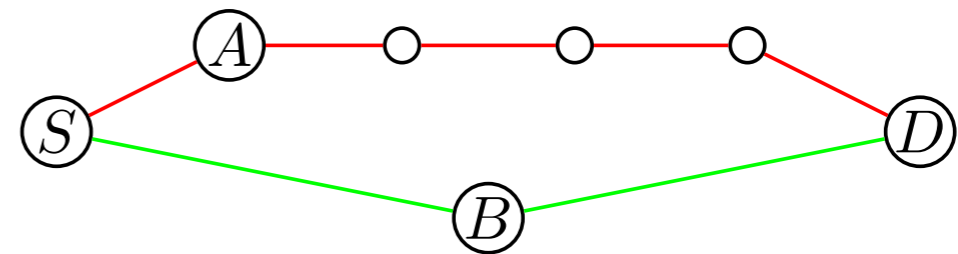
- Routing protocol for WMNs
- Ad hoc (network is not static)
- On-Demand (routes are established when needed)
- Distance (metric is hop count)
- Vector (routing table has the form of a vector)
- Developed 1997-2001 by Perkins, Beldig-Royer and Das (University of Cincinnati)

- AODV control messages
 - route request (RREQ)
 - route reply (RREP)
 - route error message (RERR)
 - (Hello messages)
- Information at nodes
 - own IP address
 - a local sequence number (freshness/timer)
 - a routing table
 - local knowledge
 - entries: (dip, dsn, val, hops, nhip, pre)

- Routing table entries (no sequence number so far)
(nhop, hops)

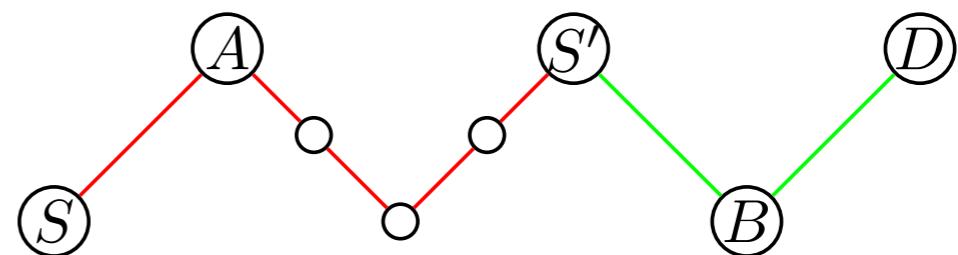
- Choice:

$$(A, 5) + (B, 2) = (B, 2)$$



- Multiplication (destination and source must coincide)

$$(A, 5) \cdot (B, 2) = (A, 7)$$



- Special symbols: $(-, 0)$, $(-, \infty)$
- Idea: back to Backhouse, Carré, Griffin, Sobrinho

- Matrices over routing table entries

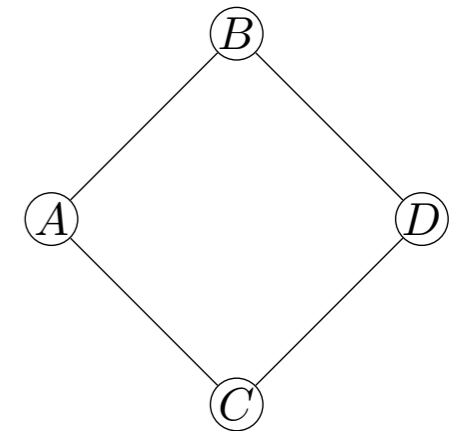
$$\begin{array}{c}
 A \\
 B \\
 C \\
 D \\
 \vdots
 \end{array}
 \begin{pmatrix}
 A & B & C & D & \dots \\
 \hline
 (-, 0) & (B, 1) & (B, 2) & (-, \infty) & \dots \\
 (A, 1) & (-, 0) & (C, 1) & (-, \infty) & \dots \\
 (-, \infty) & (B, 1) & (-, 0) & (-, \infty) & \dots \\
 (-, \infty) & (-, \infty) & (-, \infty) & (-, 0) & \dots \\
 \vdots & \vdots & \vdots & \vdots & \ddots
 \end{pmatrix}
 \begin{array}{l}
 \text{routing table of } A \\
 \\
 \\
 \\
 \\
 \end{array}$$

“routes” to B

- standard matrix operations
- further abstraction possible
(semirings, test, domain, modules ...)

Example

- A route request is broadcast



$$\begin{pmatrix} (-, 0) & (B, 1) & (C, 1) & (-, \infty) \\ (A, 1) & (-, 0) & (-, \infty) & (D, 1) \\ (A, 1) & (-, \infty) & (-, 0) & (D, 1) \\ (-, \infty) & (B, 1) & (C, 1) & (-, 0) \end{pmatrix} \cdot \begin{pmatrix} (-, 0) & (-, \infty) & (-, \infty) & (-, \infty) \\ (-, \infty) & (-, \infty) & (-, \infty) & (-, \infty) \\ (-, \infty) & (-, \infty) & (-, \infty) & (-, \infty) \\ (-, \infty) & (-, \infty) & (-, \infty) & (-, \infty) \end{pmatrix} \cdot \begin{pmatrix} (-, 0) & (B, 1) & (-, \infty) & (-, \infty) \\ (\mathbf{D}, \mathbf{3}) & (-, 0) & (-, \infty) & (-, \infty) \\ (A, 1) & (-, \infty) & (-, 0) & (D, 1) \\ (C, 2) & (-, \infty) & (C, 1) & (-, 0) \end{pmatrix}$$

topology

sender

routing table

$$= \begin{pmatrix} (-, 0) & (B, 1) & (-, \infty) & (-, \infty) \\ (\mathbf{A}, \mathbf{1}) & (-, 0) & (-, \infty) & (-, \infty) \\ (A, 1) & (-, \infty) & (-, 0) & (D, 1) \\ (C, 2) & (-, \infty) & (C, 1) & (-, 0) \end{pmatrix}$$

updated routing table

- Interpret matrix as an arbitrary element of a semiring
- Kleene algebra allows iteration,
- (Co)Domain and tests model projections

- sending messages

$$a + p \cdot b \cdot q \cdot (1 + c)$$

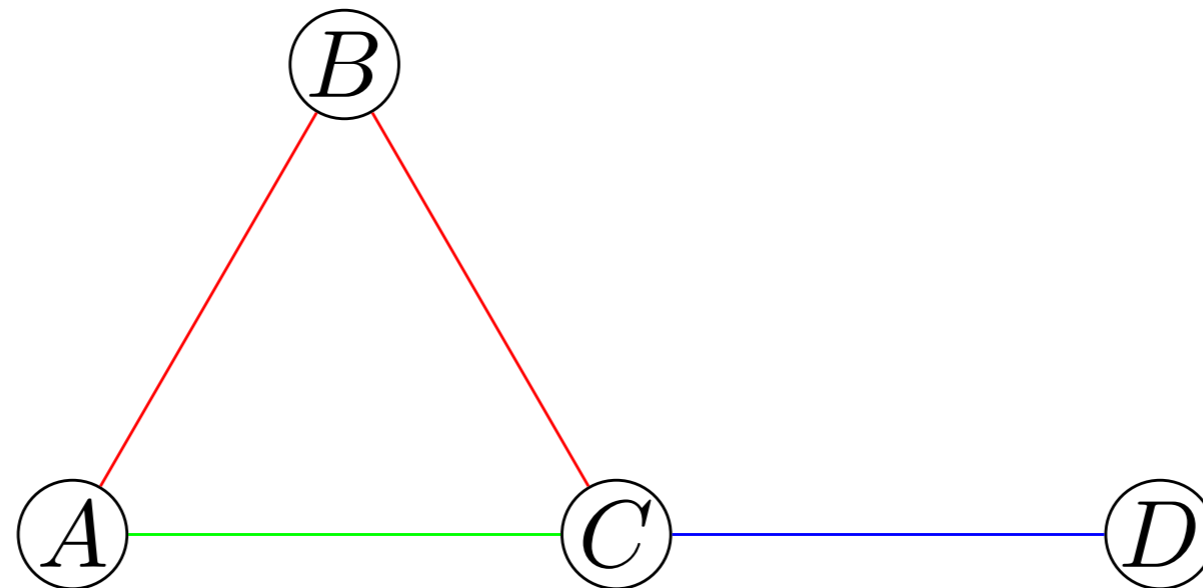
- by distributivity

$$a + p \cdot b \cdot q + p \cdot b \cdot q \cdot c$$

snapshot, 1-hop connection learnt, content sent

- broadcast, unicast, groupcast are the same (modelled by different topologies)
- Kleene star models flooding the network (modal operators terminate flooding)
- QUESTION: Can unicast modelled purely algebraically?

- Adding sequence numbers



$$r \cdot b = (B, 2, 5) \cdot (D, 1, 10) = (B \cdot D, 2 + 1, \max(5, 10)) = (B, 3, 10)$$

$$g \cdot b = (C, 1, 3) \cdot (D, 1, 10) = (C \cdot D, 1 + 1, \max(3, 10)) = (C, 2, 10)$$

$$r \cdot b + g \cdot b \neq (r + g) \cdot b$$

- **Restrict multiplication**
 - partial defined operation
 - only topologies allowed on the left-hand side
 - Kleene star has to be adapted
- **Module like structure**
(scalars are subalgebra)

- Ad hoc prototype in Haskell
- Theorems at algebraic level proven with Prover9
- Use Isabelle/HOL to switch between model and algebra

- Include sequence numbers
- Important properties loop freedom, route correctness
- Improvement/refinement
- So far concentrated on AODV
 - well known, IETF standard, known limitations
- Extend formal methods to other protocols
 - OSLR, DYMO, DLR ...
- Add further necessary concepts
 - time, probability

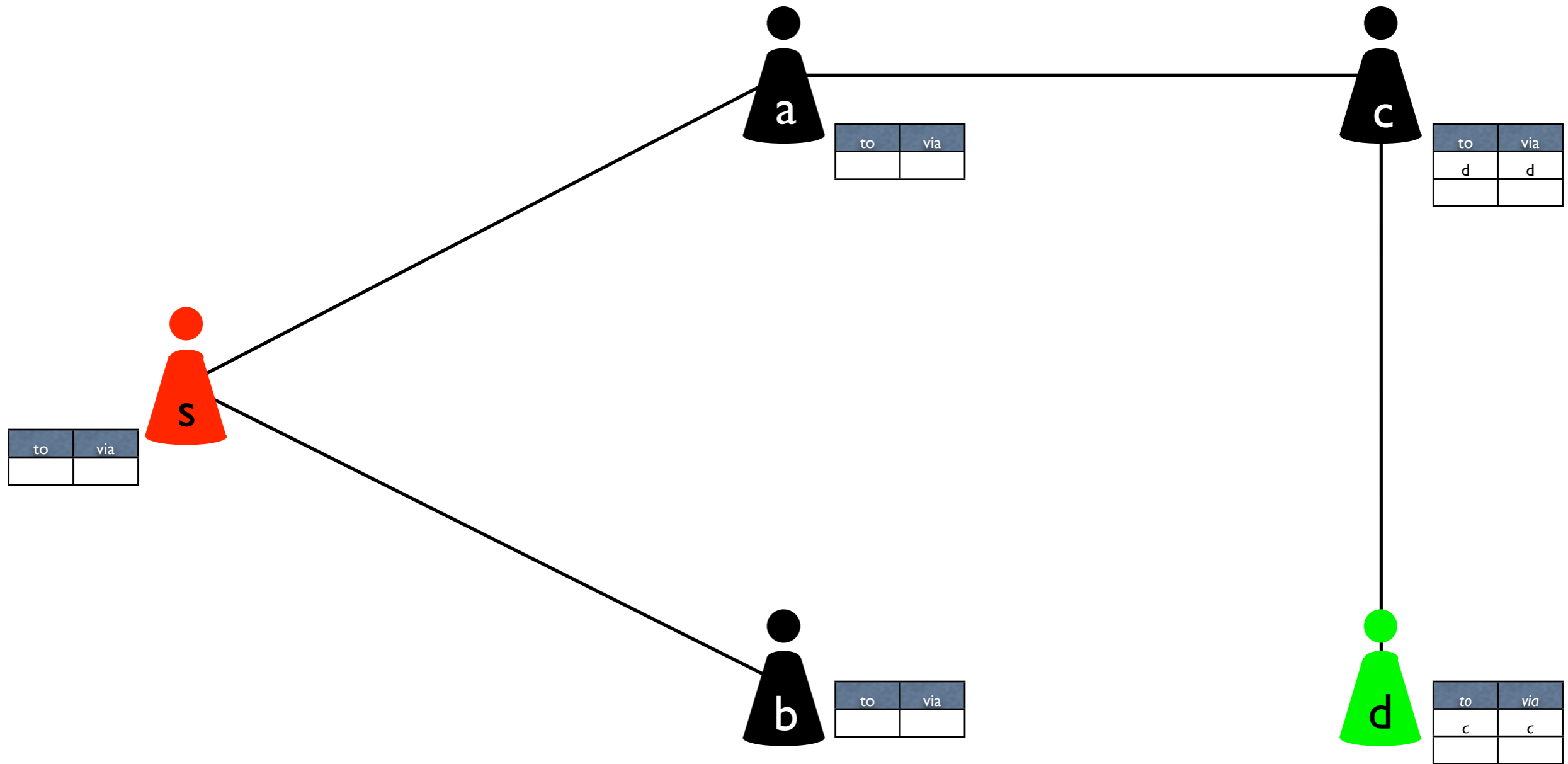


From imagination to **impact**



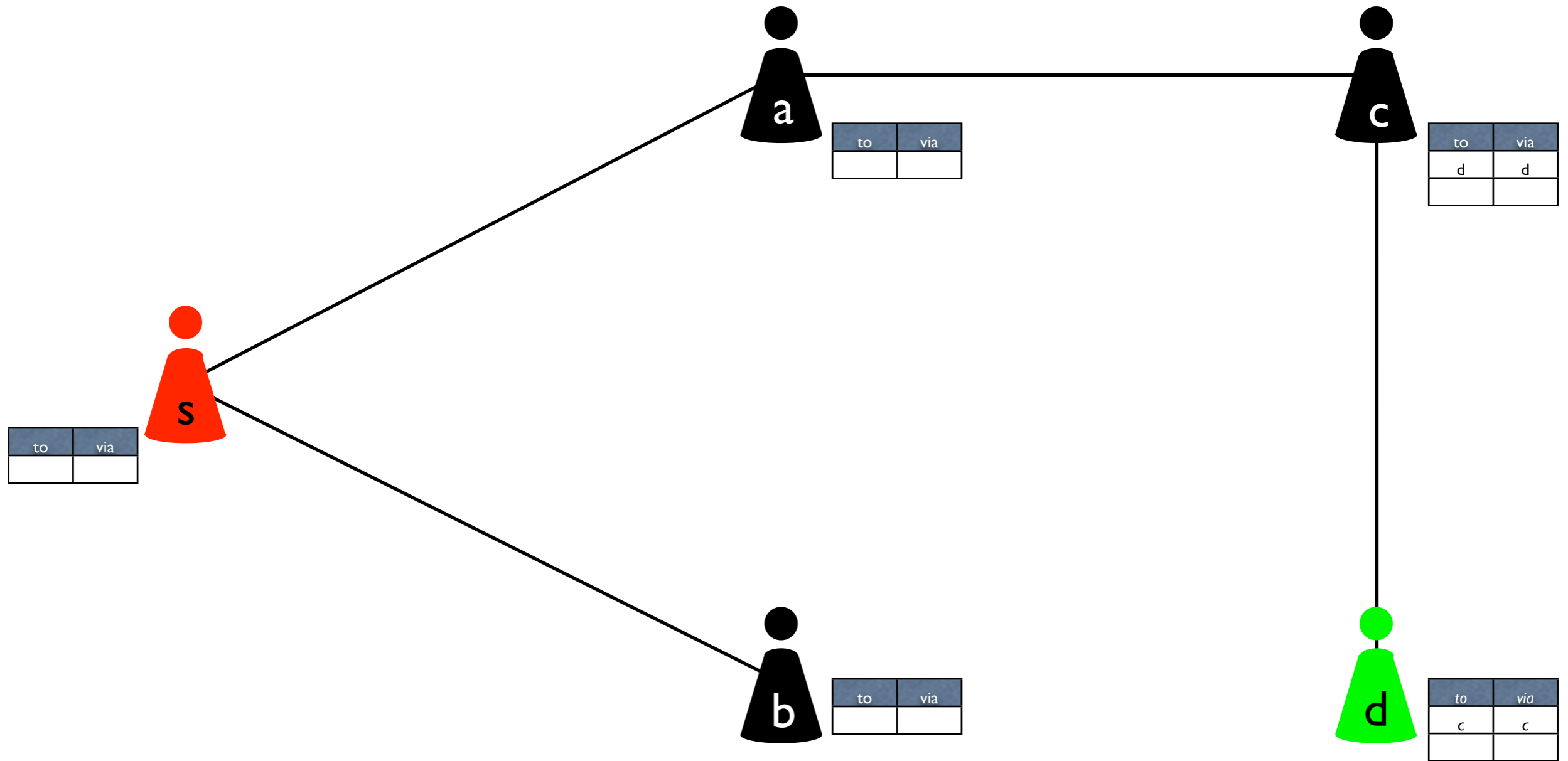
From imagination to **impact**

AODV – An Example

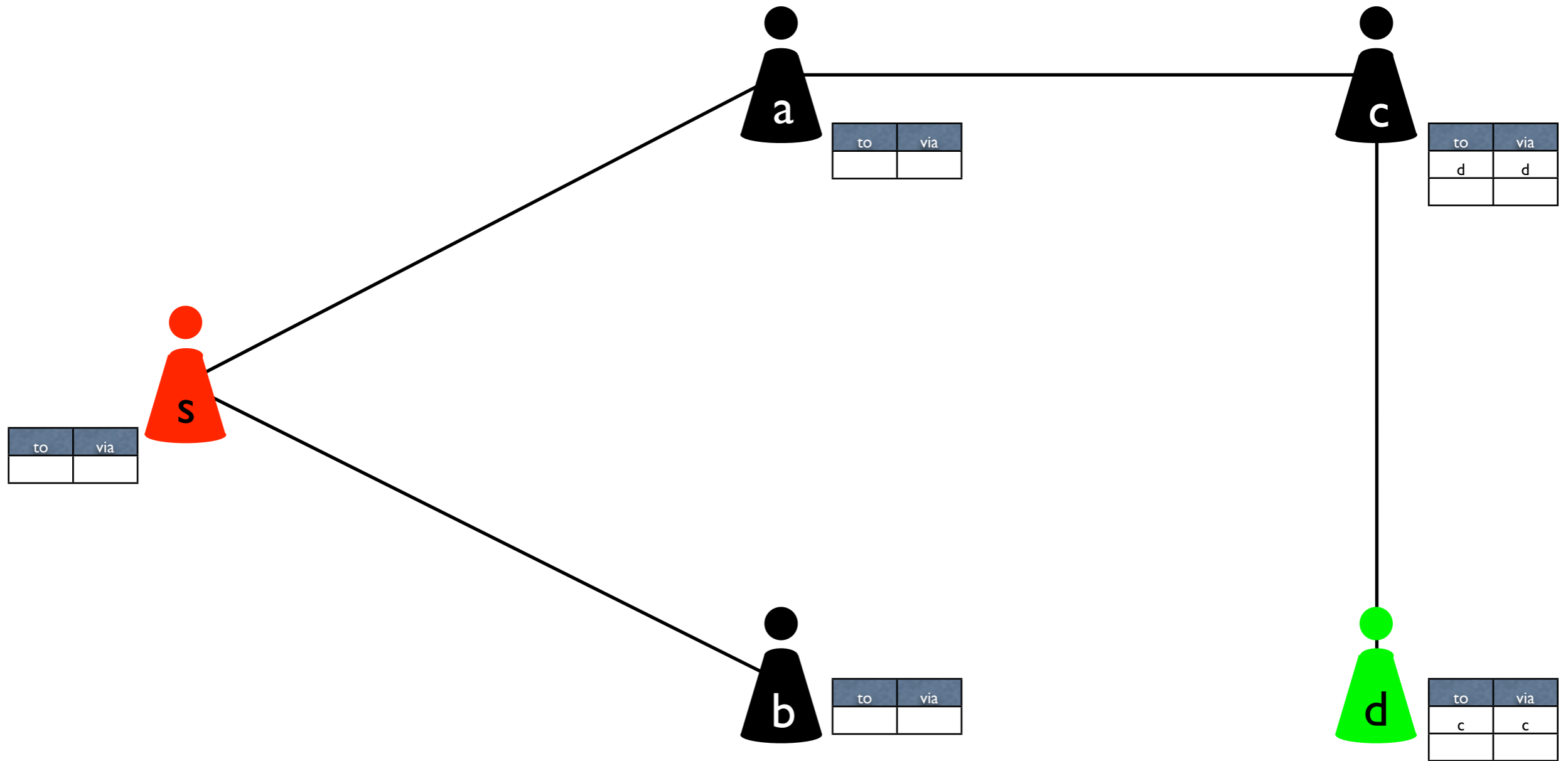


s is looking for a route to d

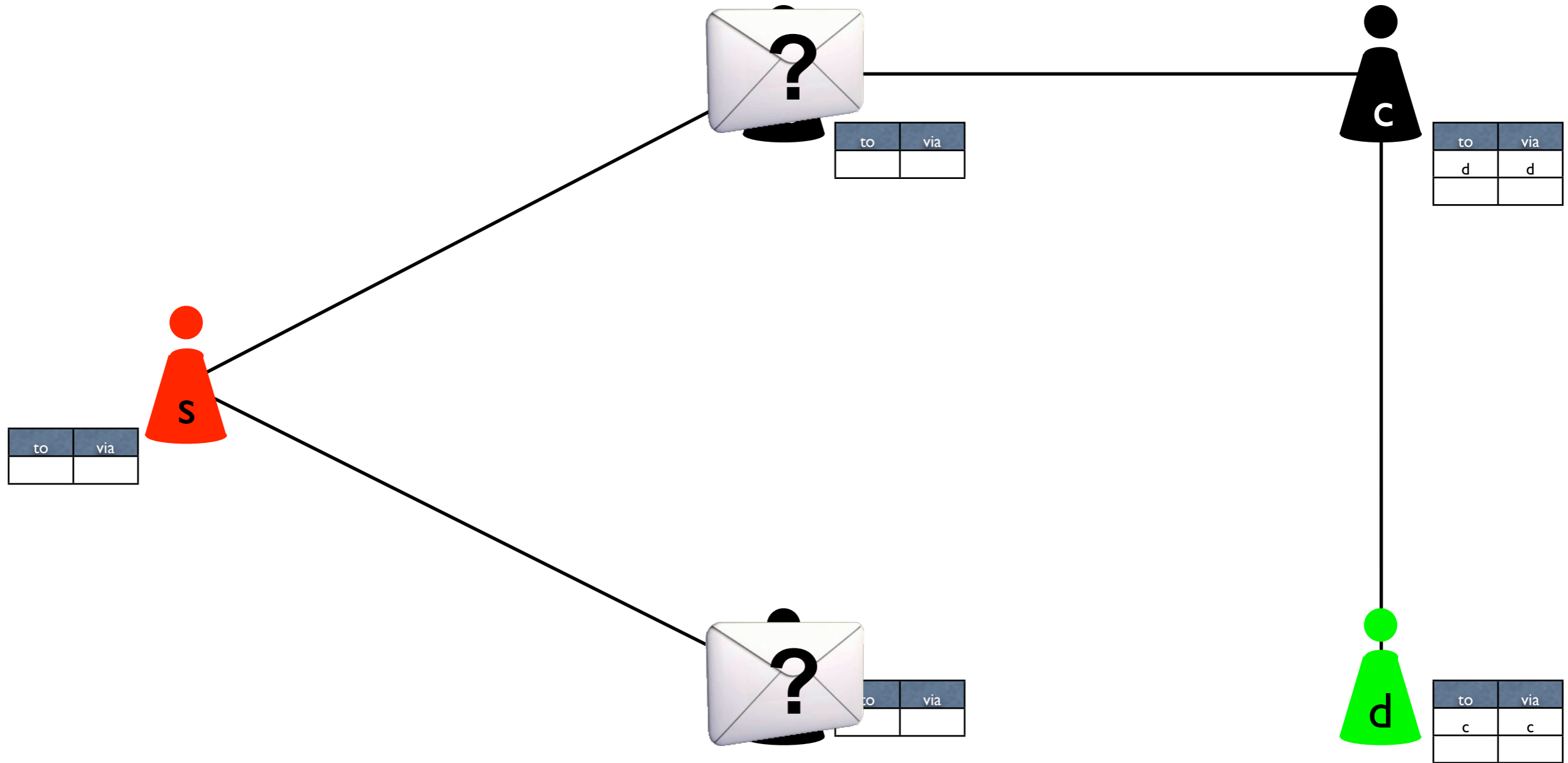
AODV – An Example



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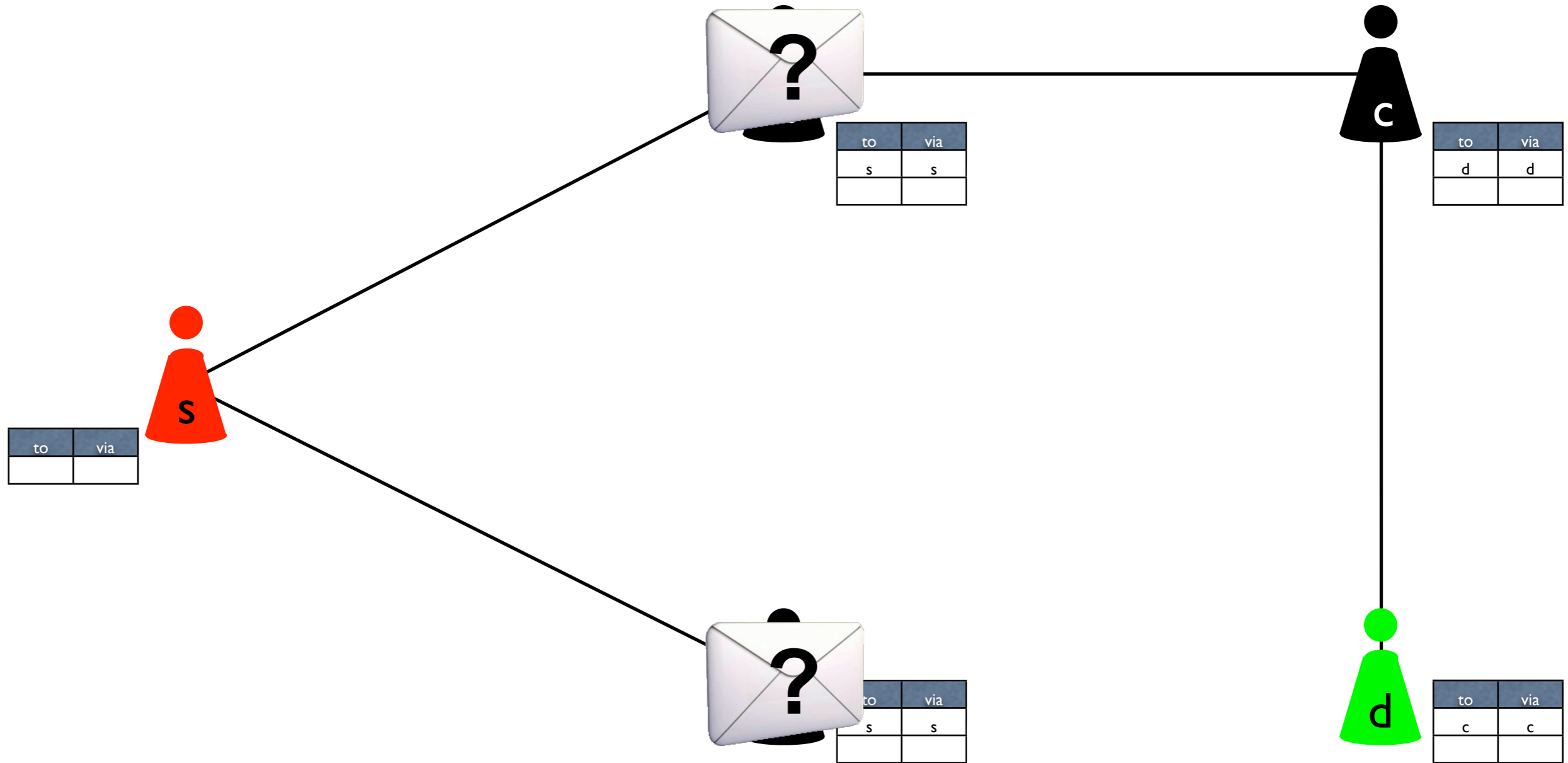


AODV – An Example



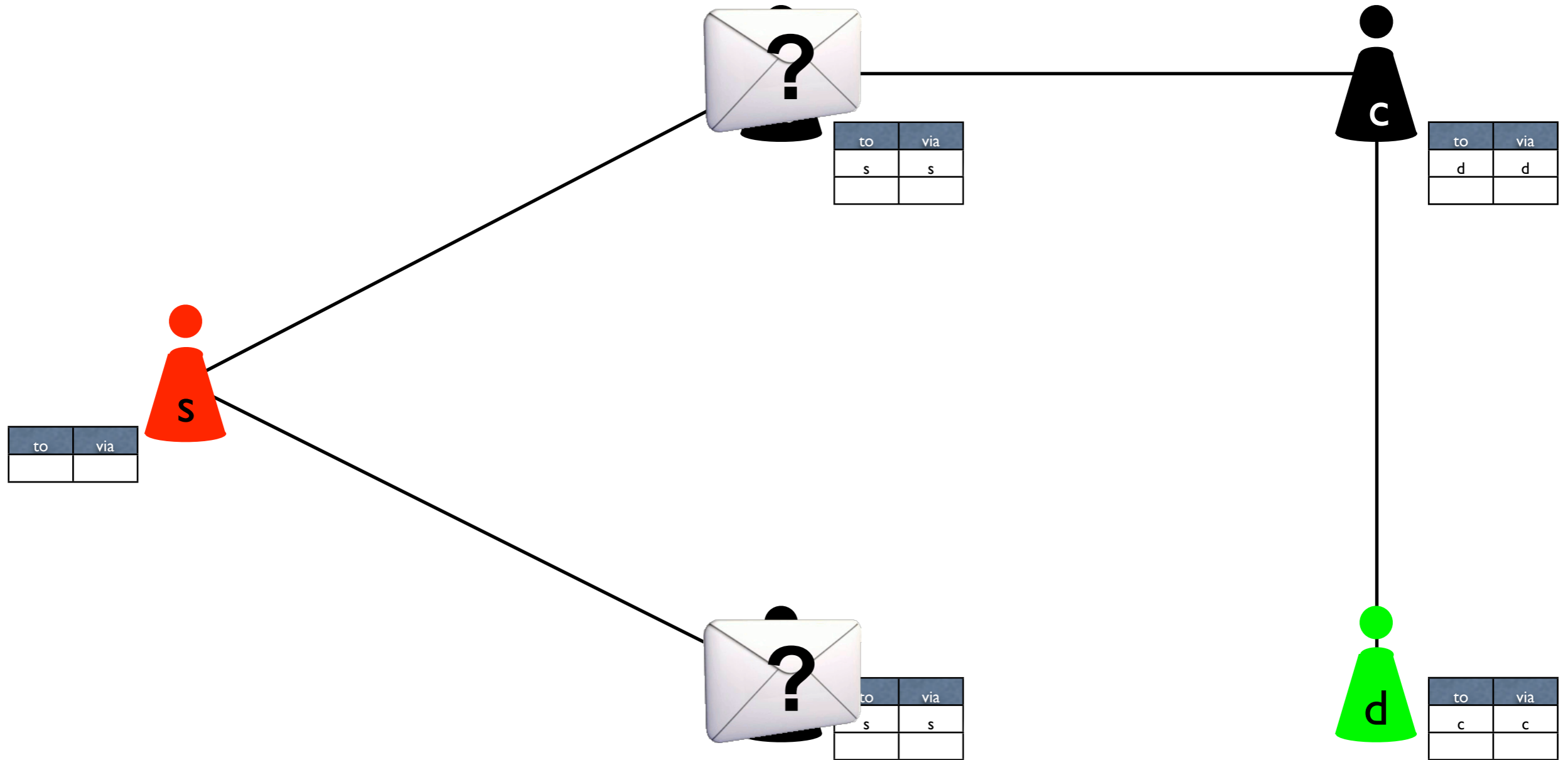
s broadcasts a route request

AODV – An Example

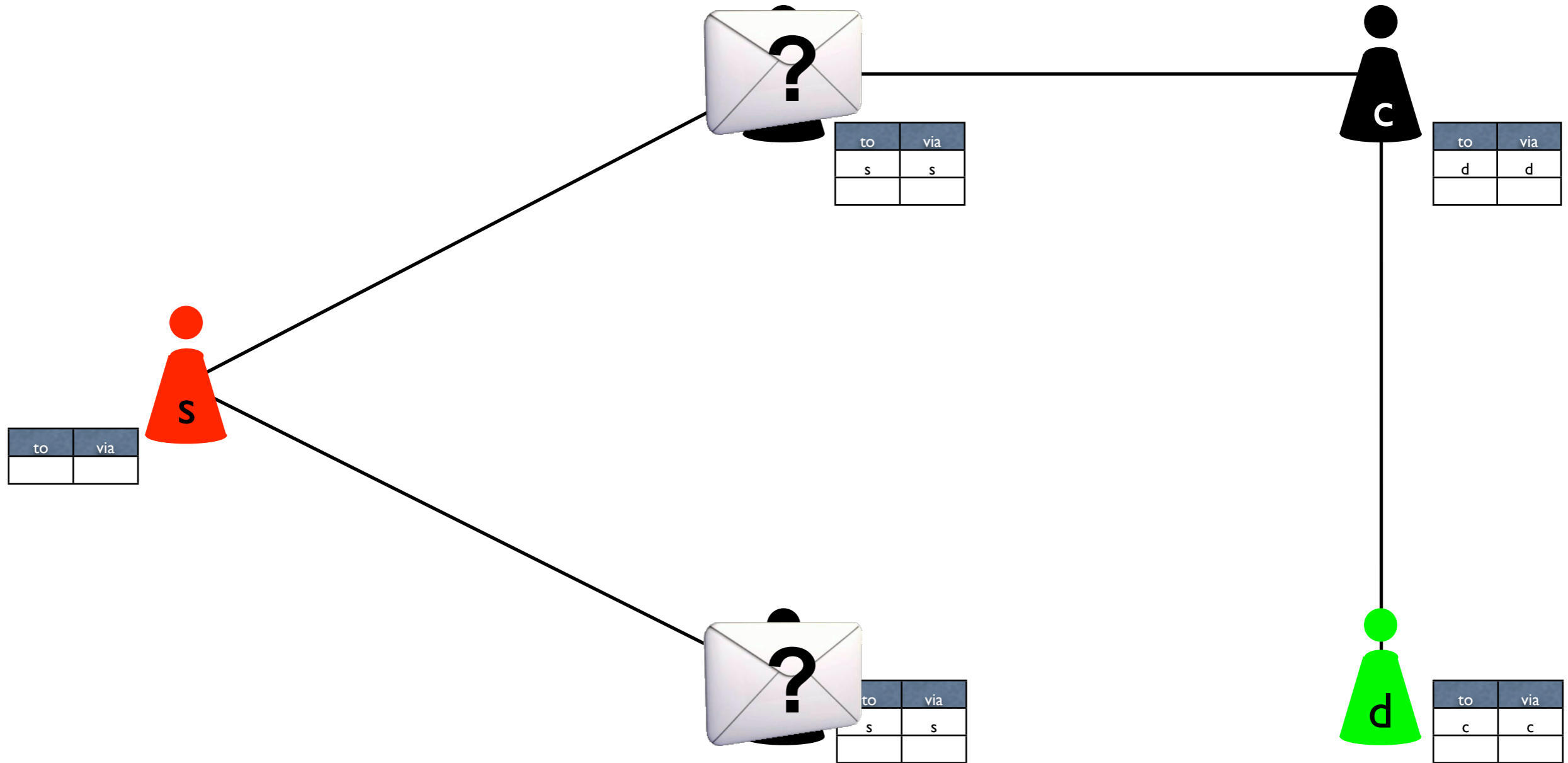


s broadcasts a route request

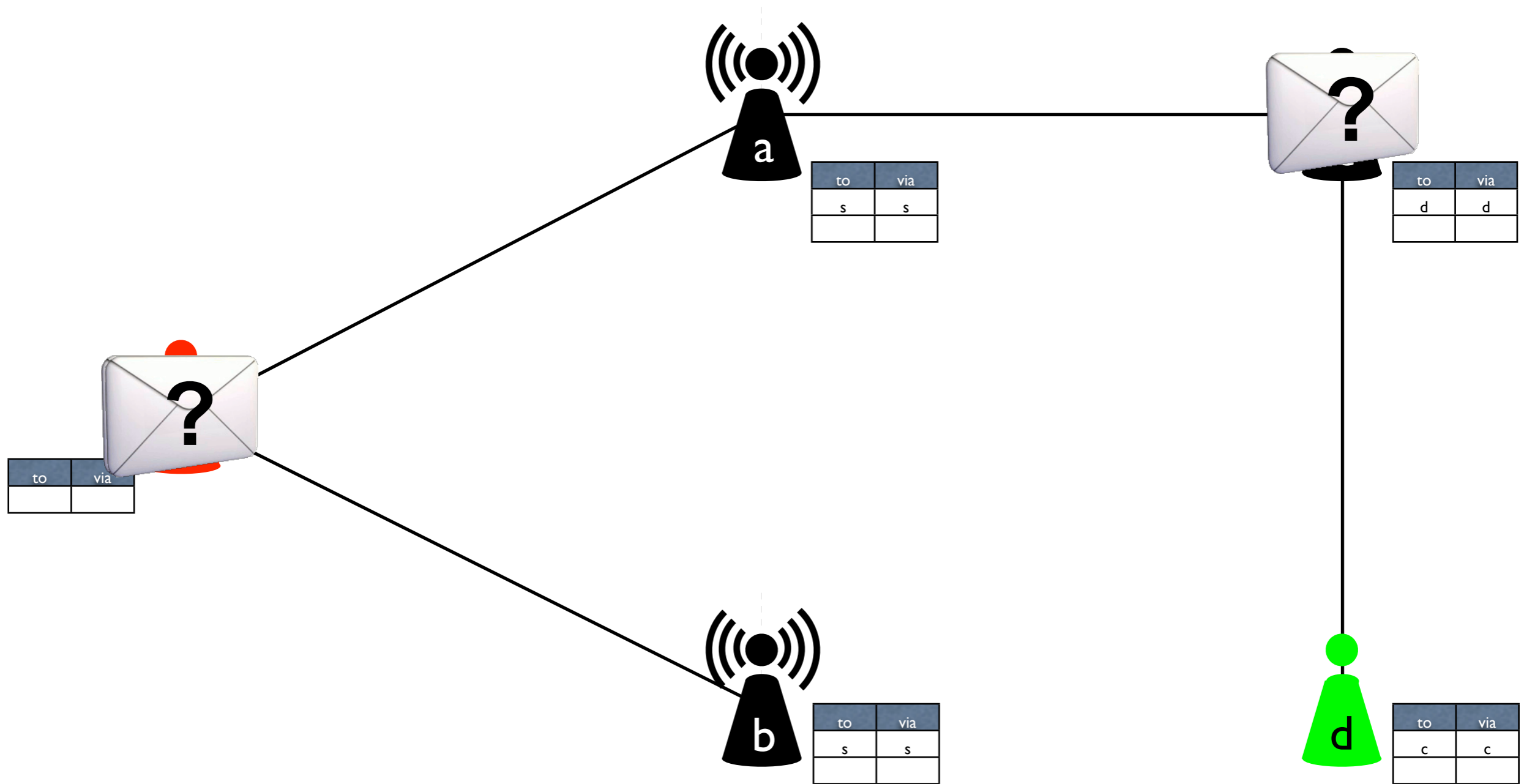
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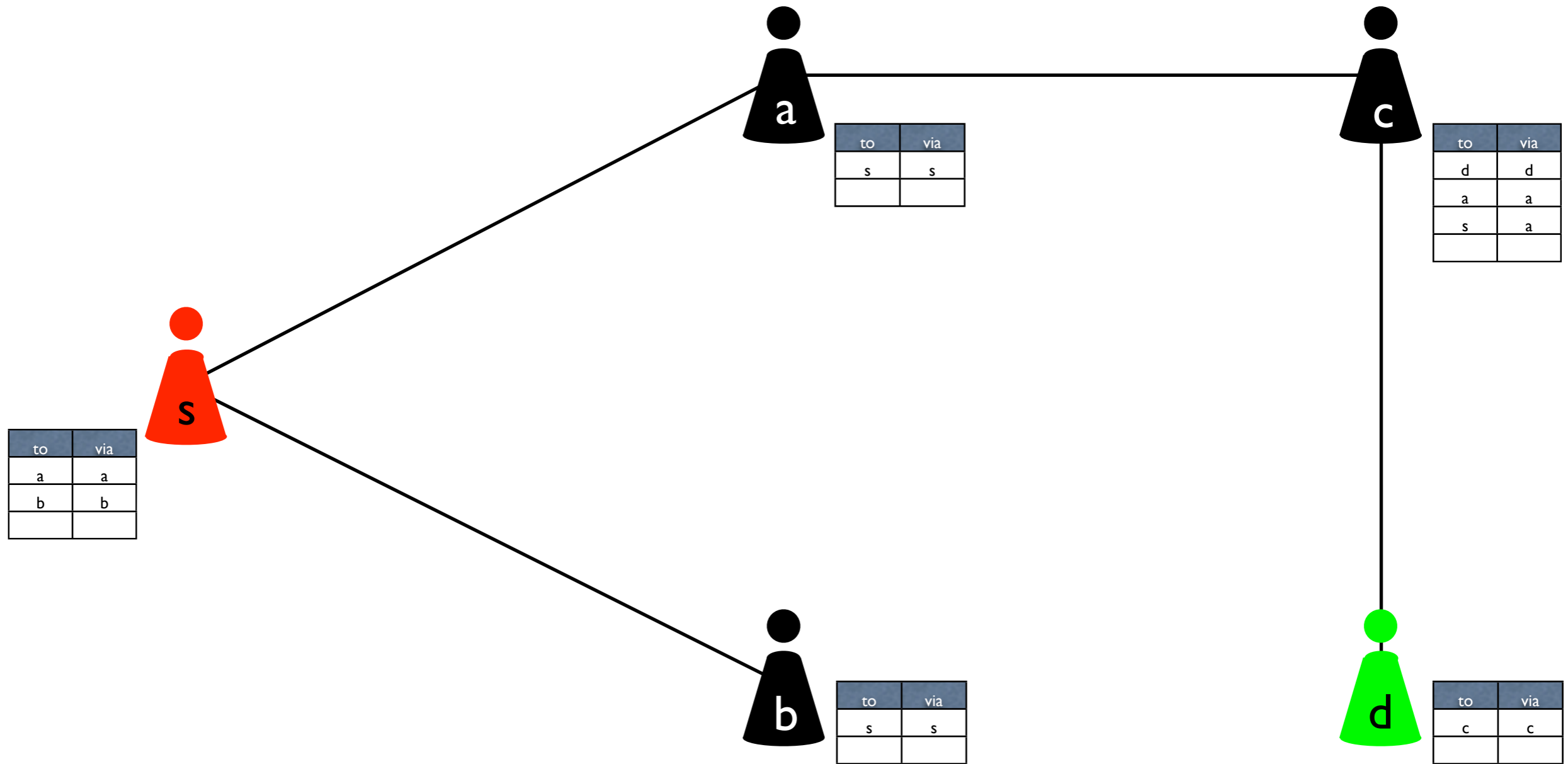


AODV – An Example



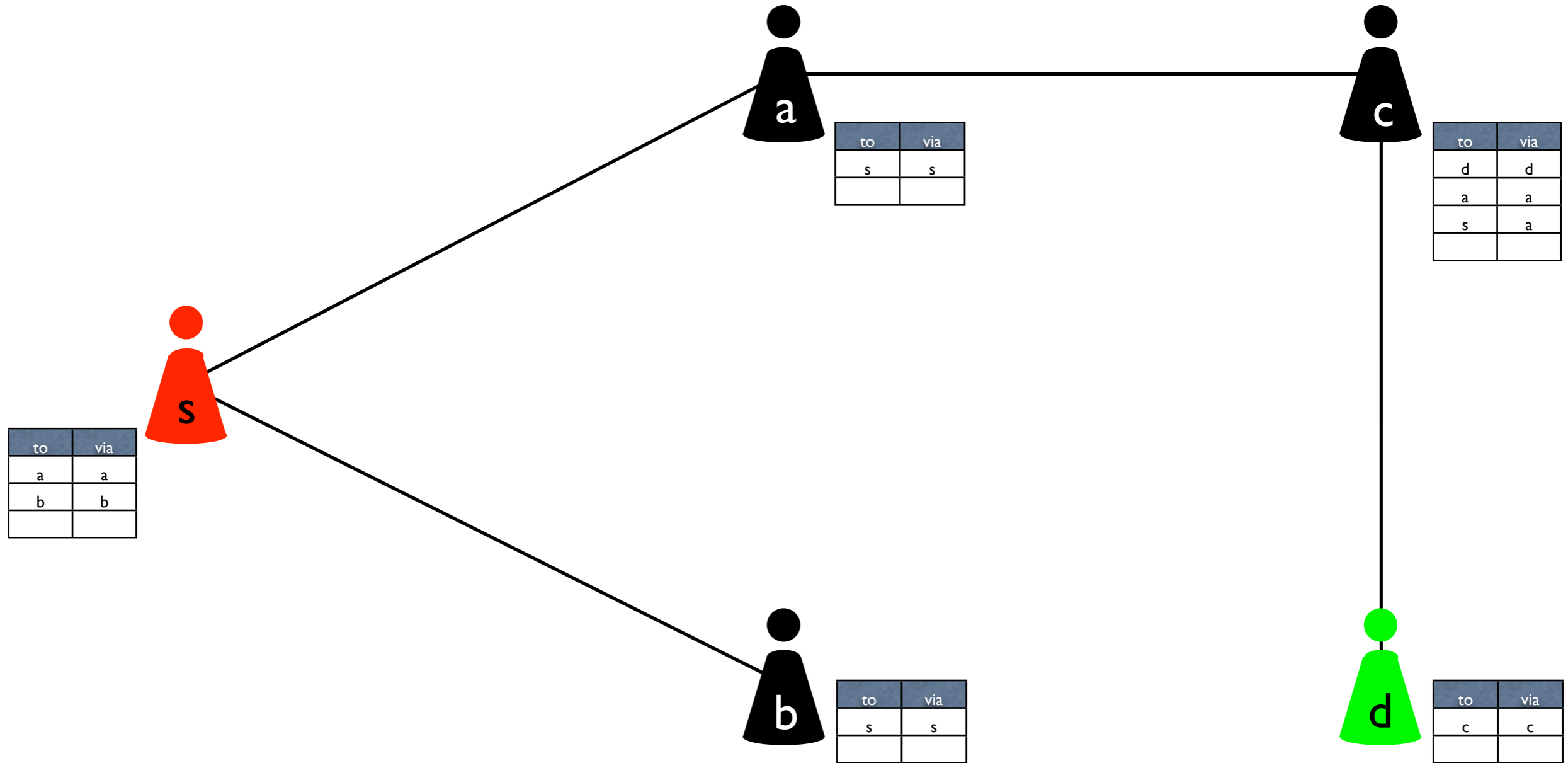
a,b forward the route request

AODV – An Example

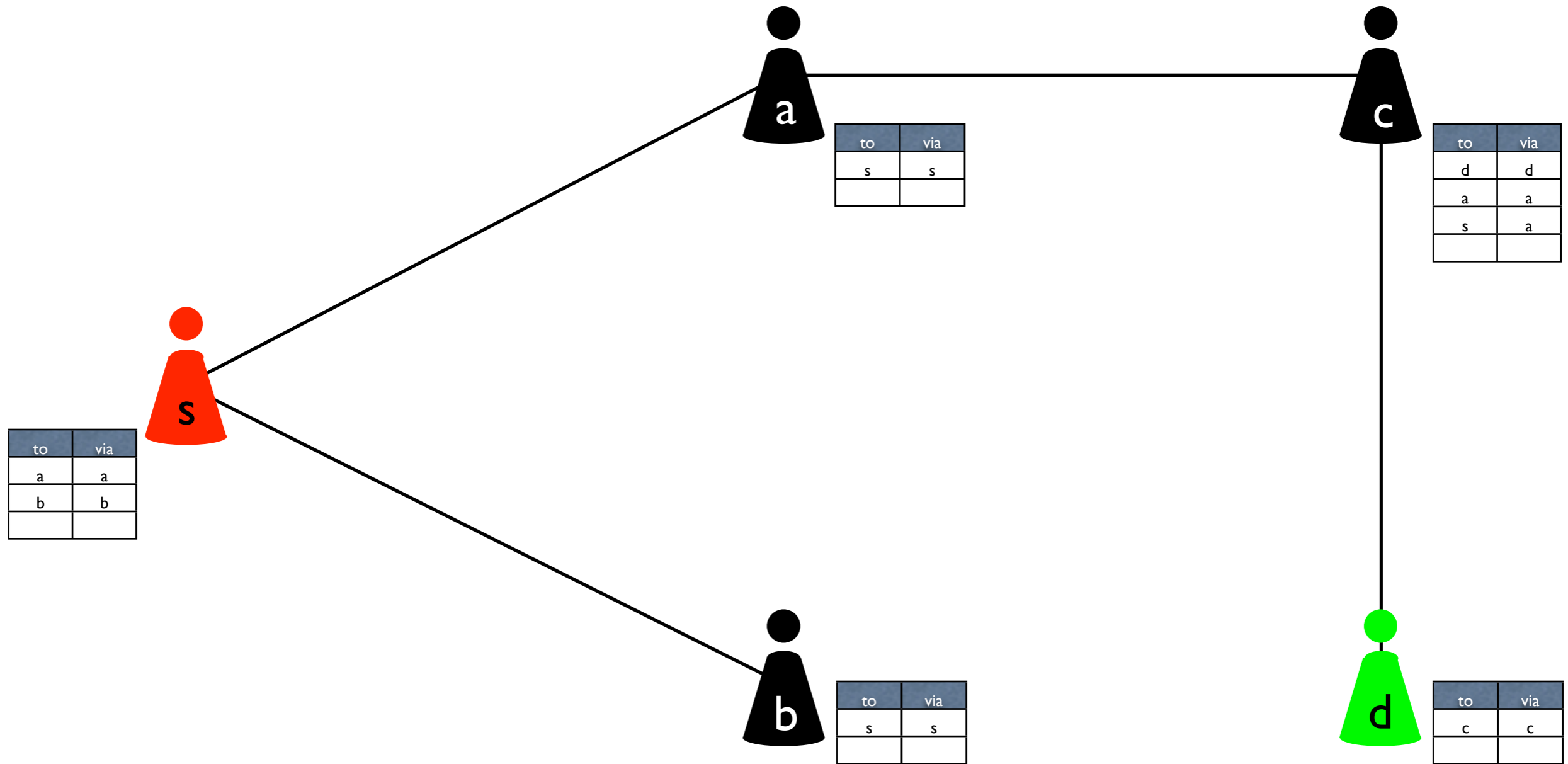


a,b forward the route request

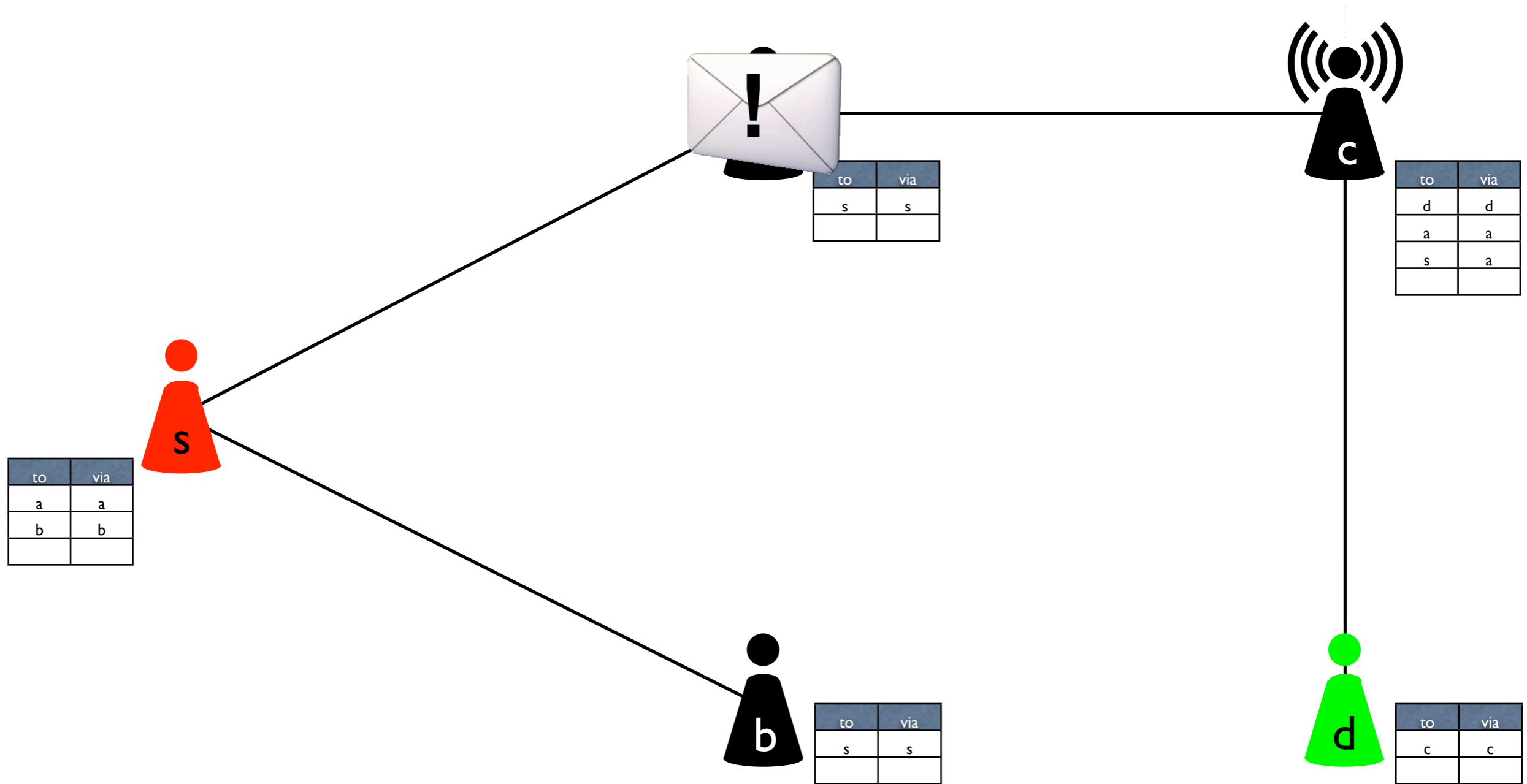
AODV – An Example



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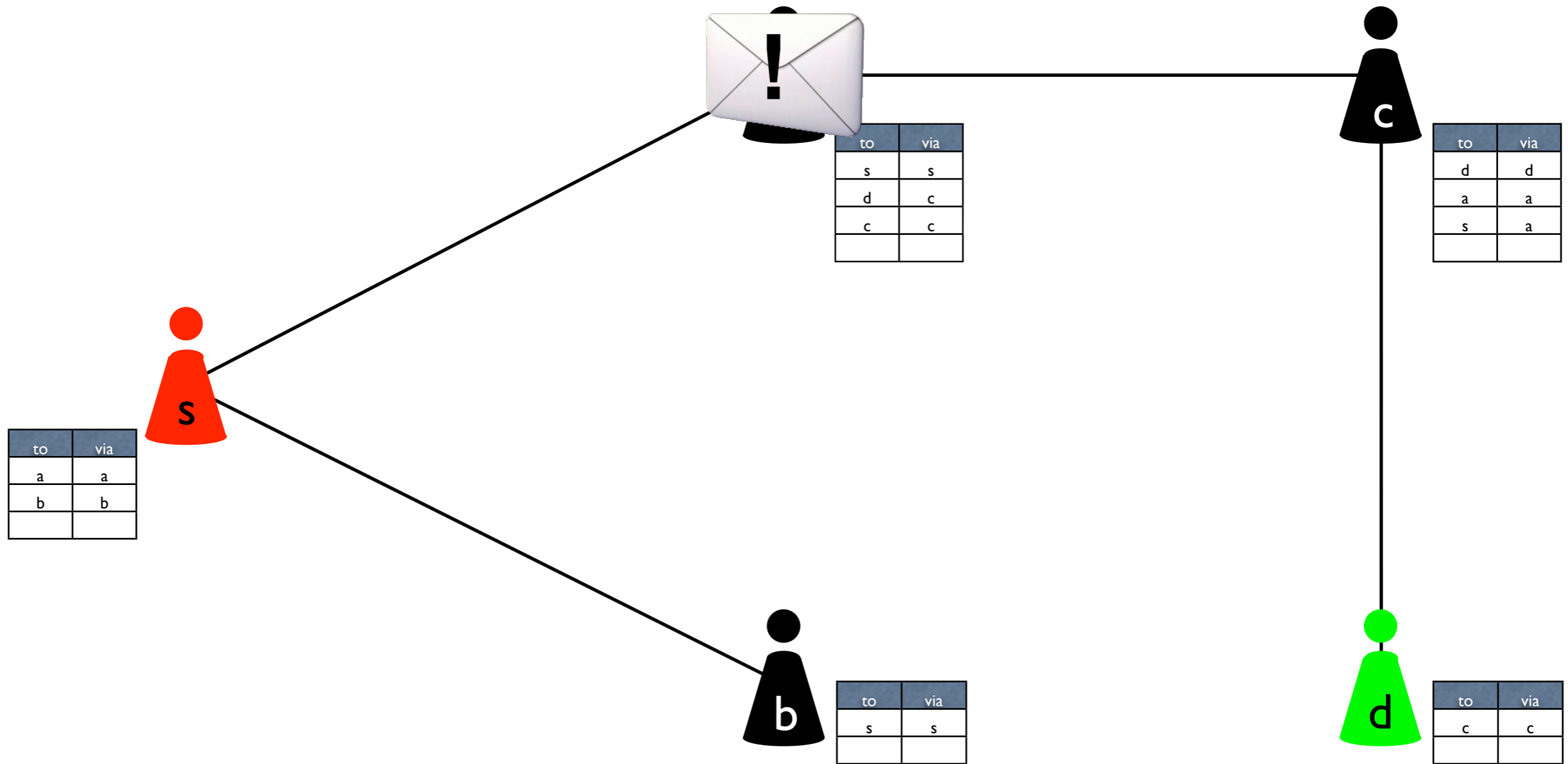


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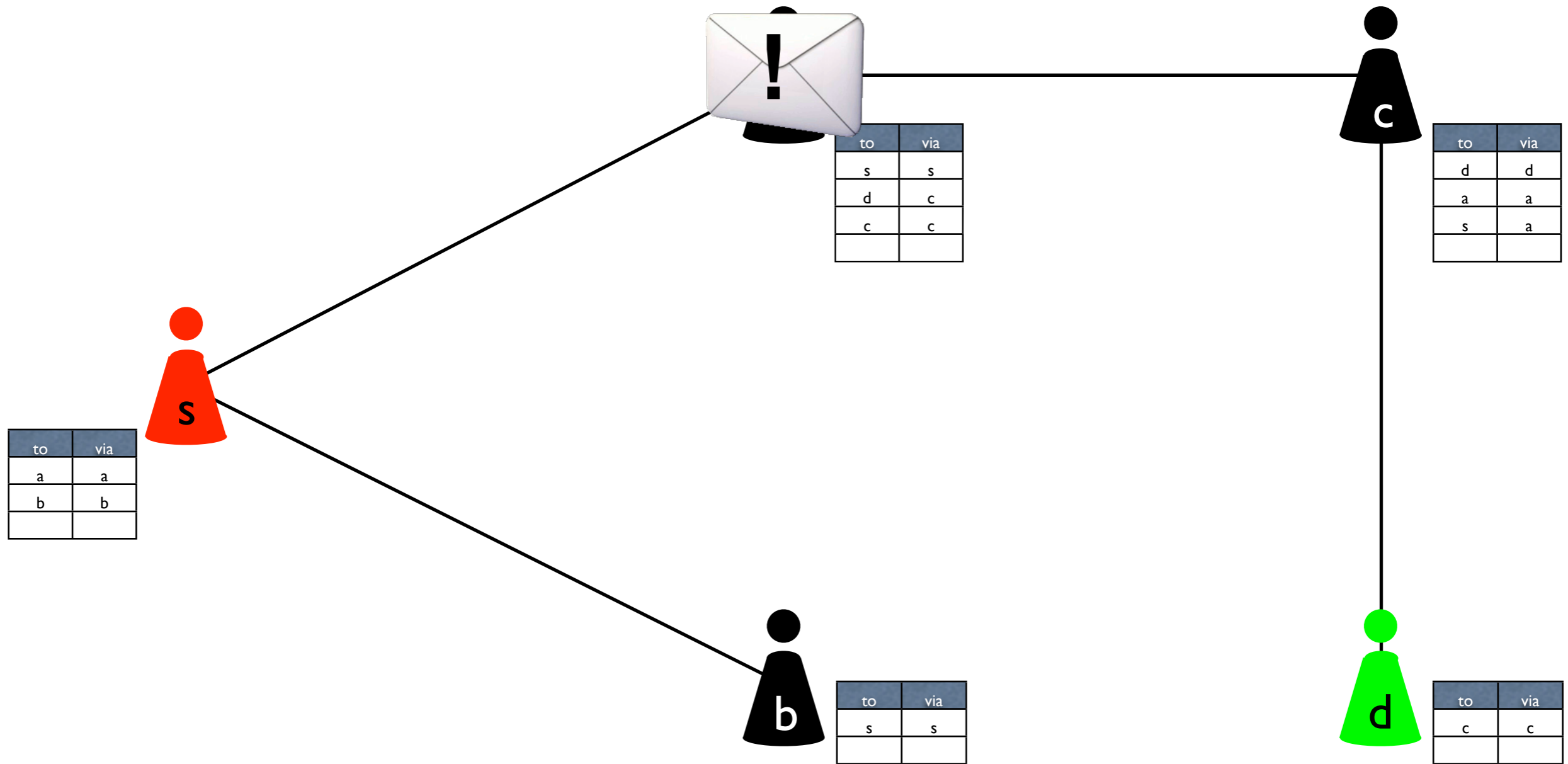
c has information about d
c answers route request and sends reply

AODV – An Example

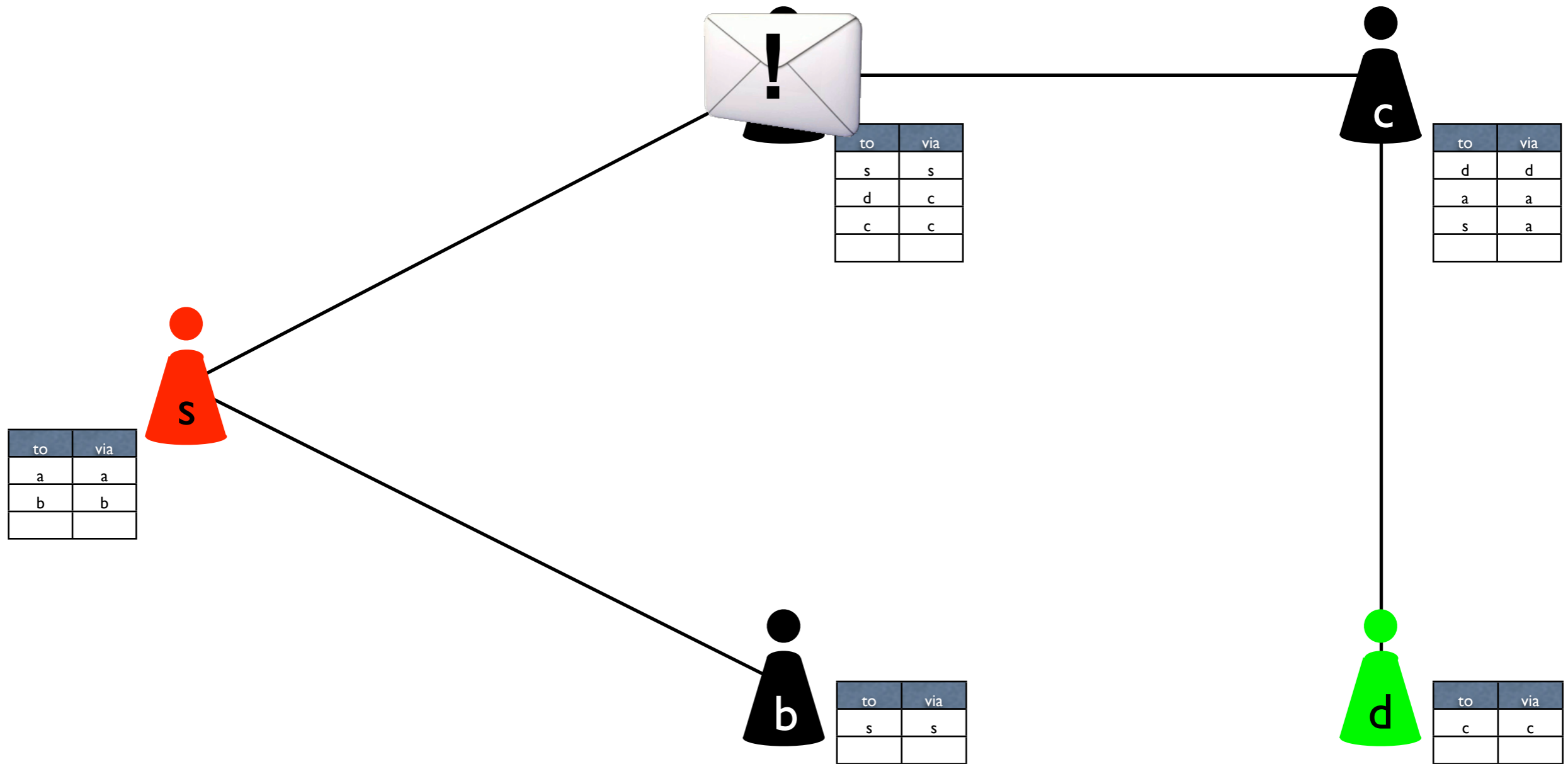


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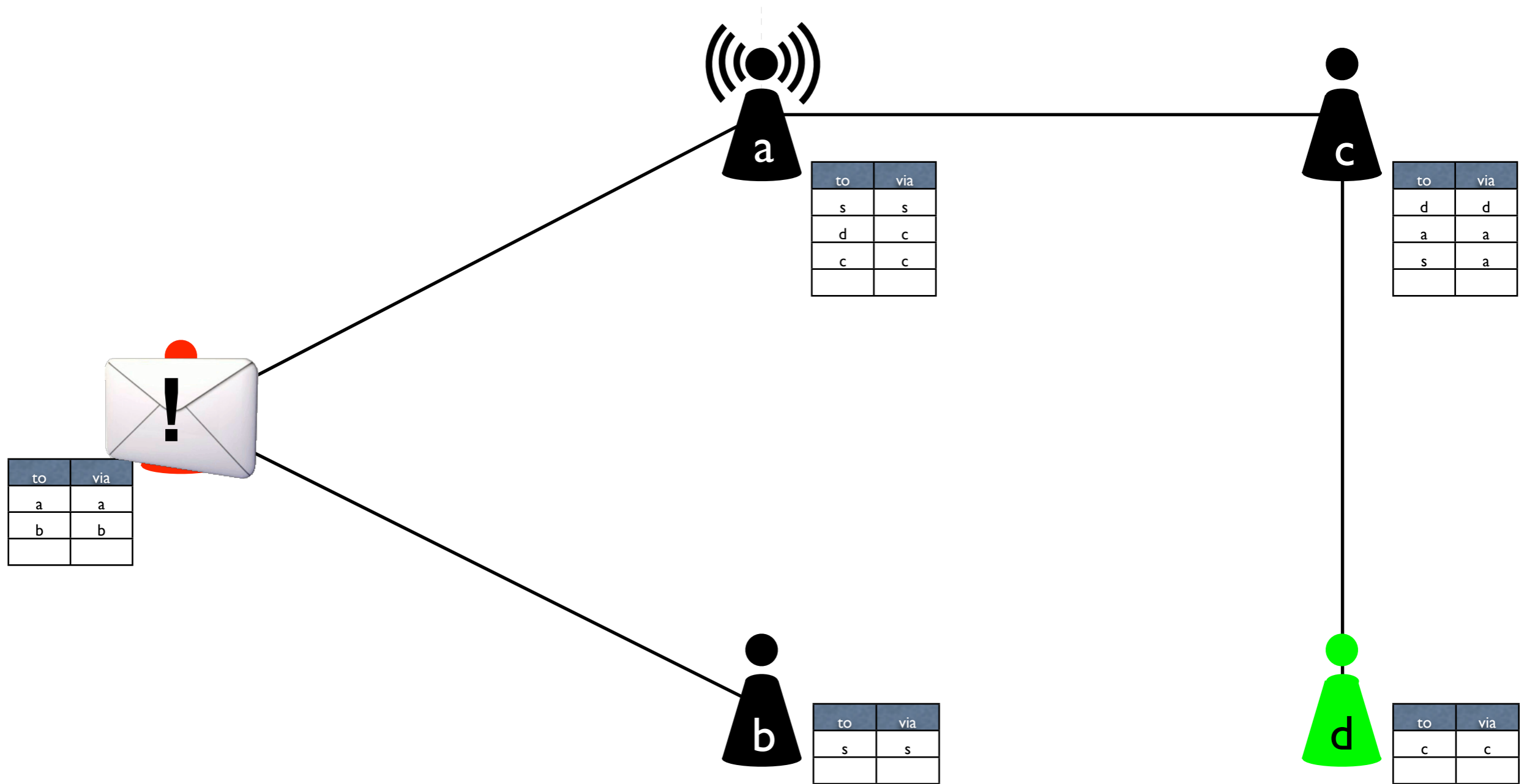
AODV – An Example



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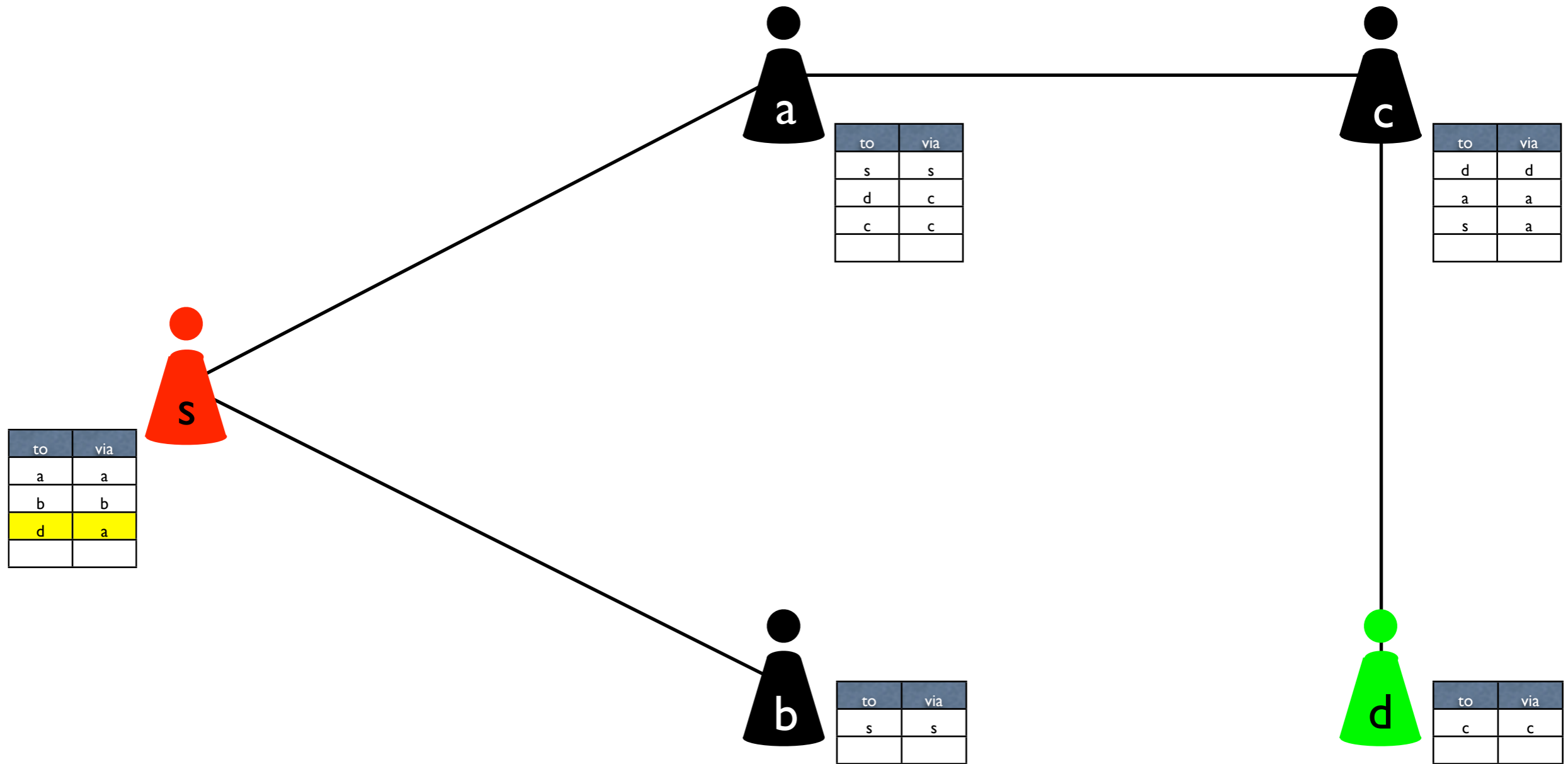


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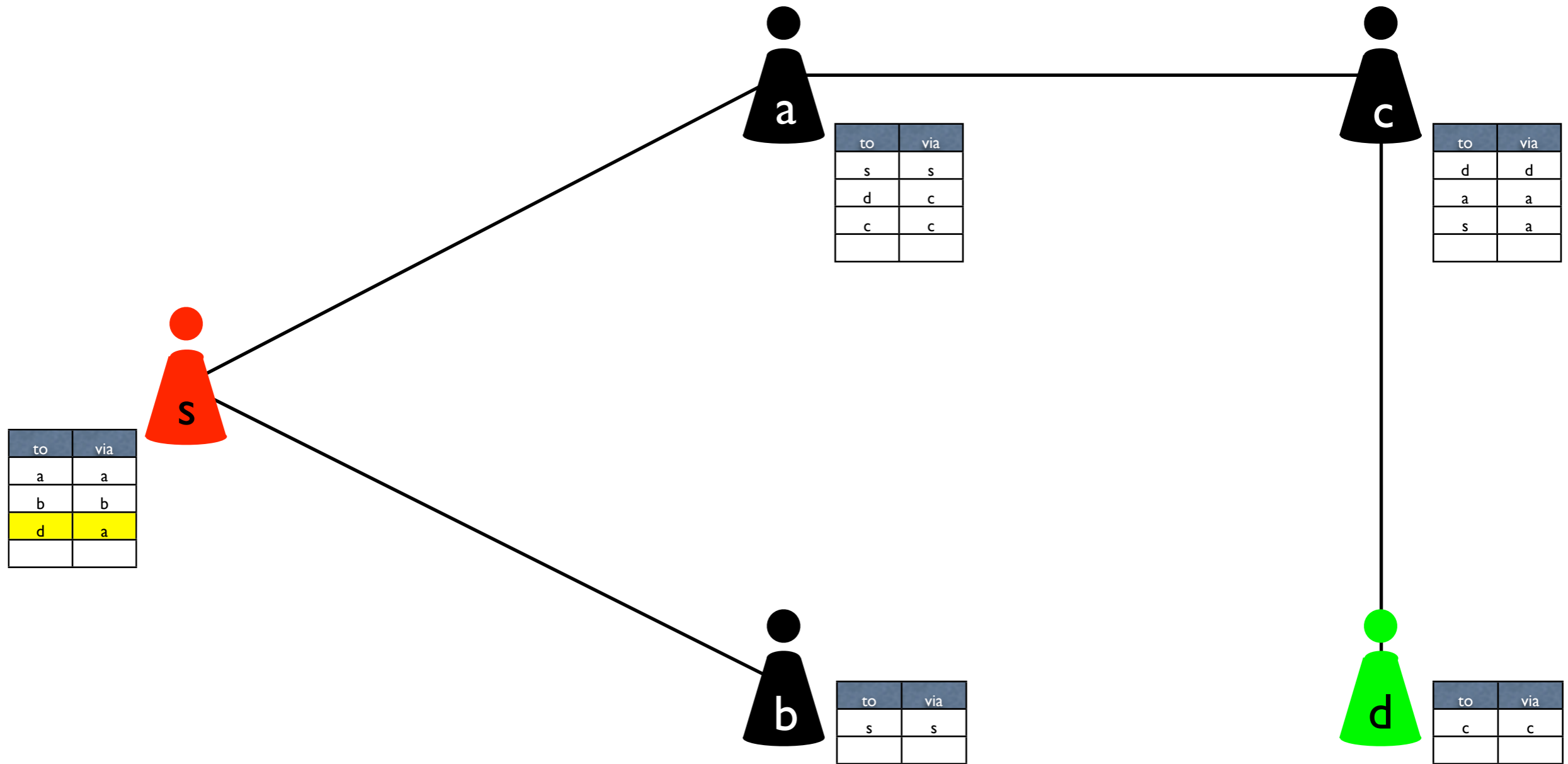
a forwards route reply

AODV – An Example

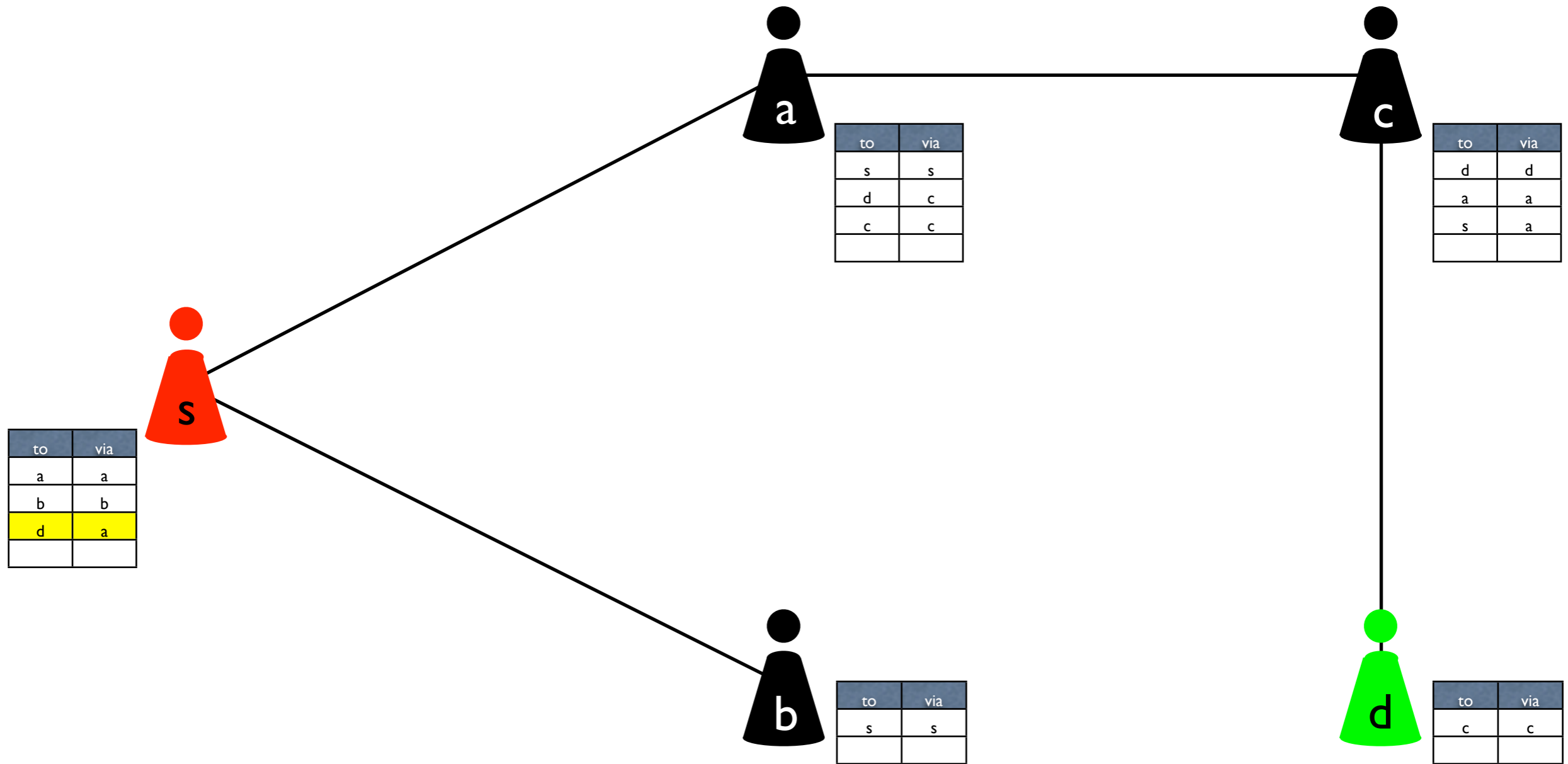


a forwards route reply

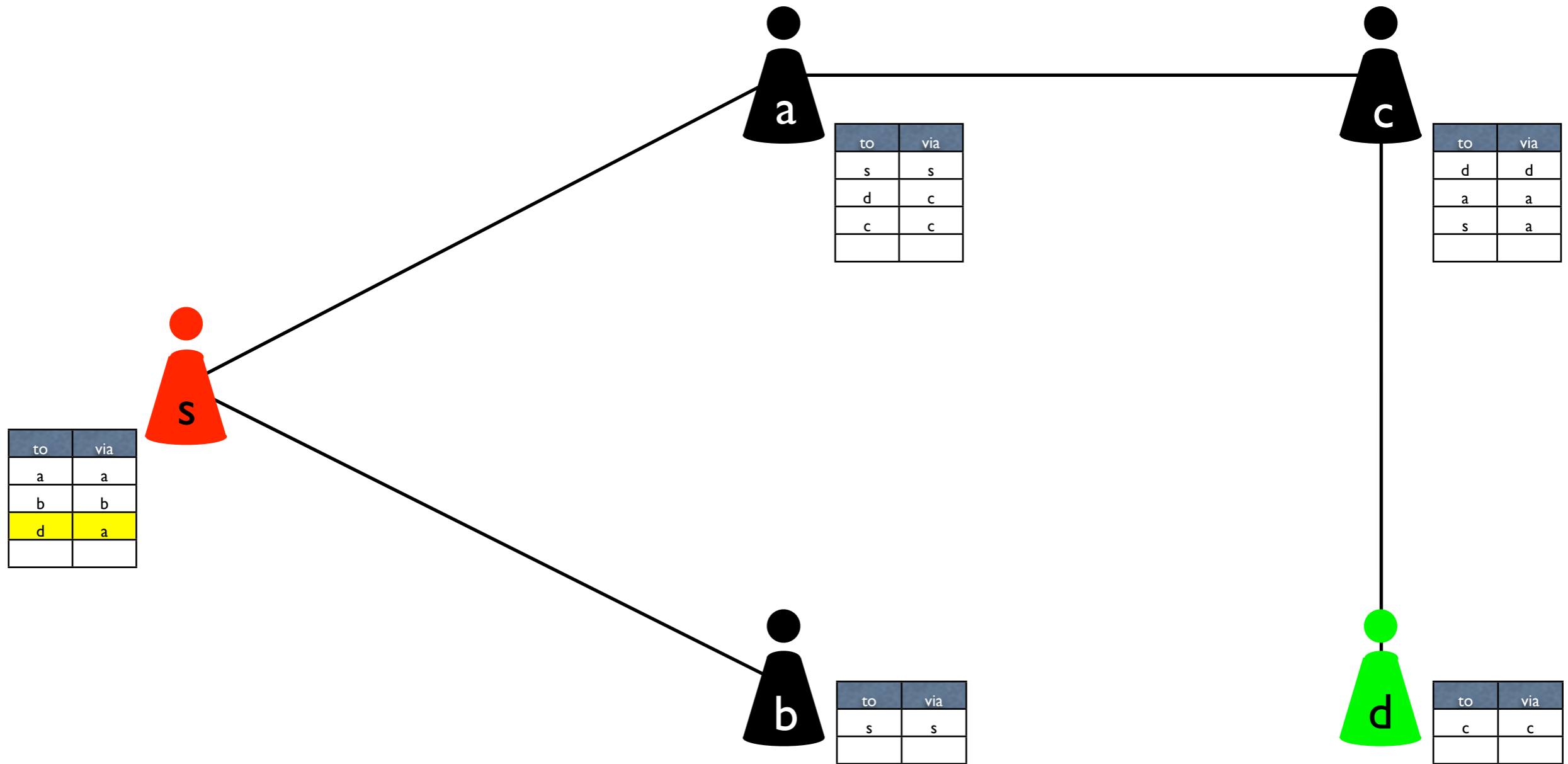
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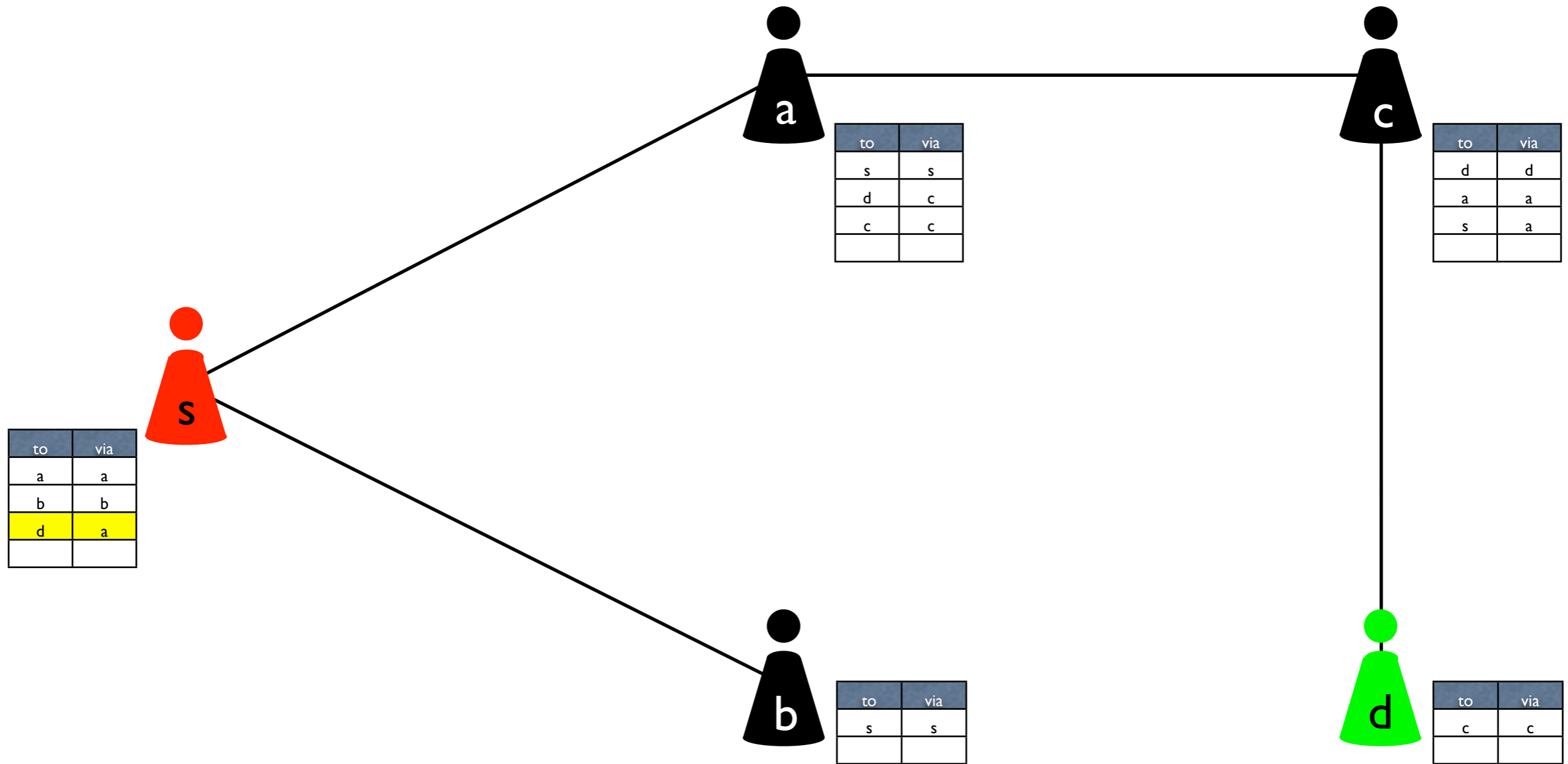


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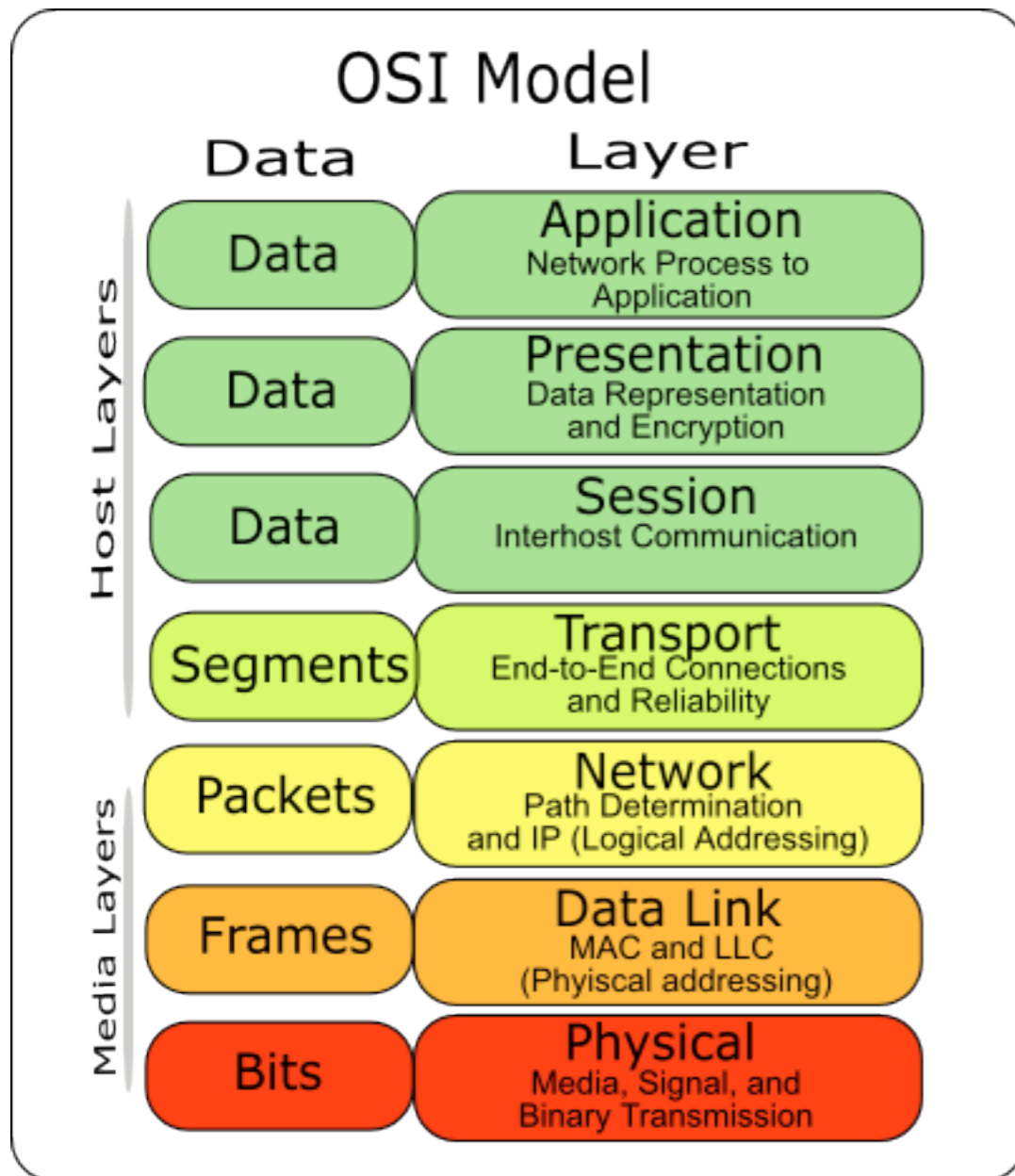
s has found a route to d

AODV – An Example



s has found a route to d

Different Network Layers



- Routing protocols
 - find (optimal) route
 - properties
 - loop freedom (no packet travels in loops)
 - route correctness (if a route is found, the route is valid)
 - route found (if a route exists, at least one route is found)
 - packet delivery
- Routing tables
 - data structure
 - belongs to client/router
 - lists destinations
 - sometimes metrics